

# NetSMF: Large-Scale Network Embedding as Sparse Matrix Factorization

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# Motivation and Problem Formulation

## Problem Formulation

Give a network  $G = (V; E)$ , aim to learn a function  $f: V \rightarrow \mathbb{R}^p$  to capture neighborhood similarity and community membership.

## Applications:

- | link prediction
- | community detection
- | label classification

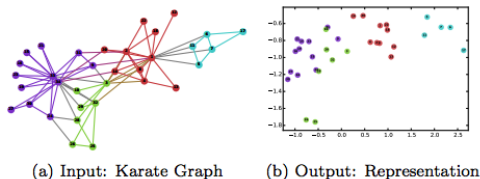


Figure 1: A toy example (Figure from DeepWalk).

# Two Genres of Network Embedding Algorithm

- | Local Context Methods:
  - ▶ LINE, DeepWalk, node2vec, metapath2vec.
  - ▶ Usually be formulated as a skip-gram-like problem, and optimized by SGD.
- | Global Matrix Factorization Methods.
  - ▶ NetMF, GraRep, HOPE.
  - ▶ Leverage global statistics of the input networks.
  - ▶ Not necessarily a gradient-based optimization problem.
  - ▶ Usually requires explicit construction of the matrix to be factorized.

# Notations

Consider an undirected weighted graph  $G = (V; E)$ , where  $|V| = n$  and  $|E| = m$ .

- | Adjacency matrix  $\mathbf{A} \in \mathbb{R}_+^{n \times n}$ :

$$\mathbf{A}_{i,j} = \begin{cases} a_{i,j} > 0 & (i,j) \in E \\ 0 & (i,j) \notin E \end{cases}$$

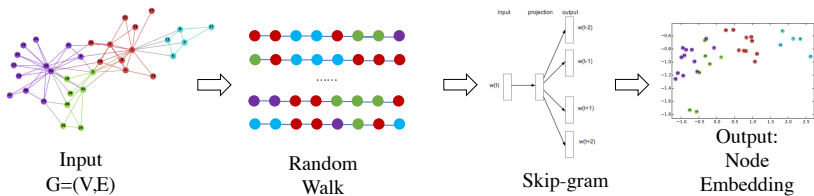
- | Degree matrix  $\mathbf{D} = \text{diag}(d_1; \dots; d_n)$ , where  $d_i$  is the generalized degree of vertex  $i$ .
- | Volume of the graph  $G$ :  $\text{vol}(G) = \sum_i \sum_j \mathbf{A}_{i,j}$ .

Revisit DeepWalk and NetMF

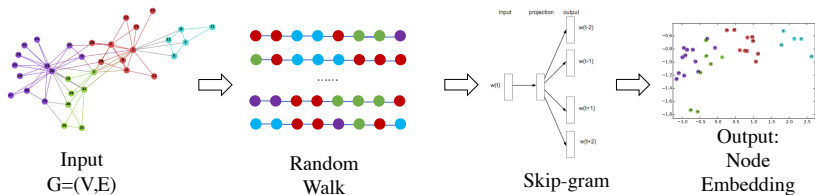
NetSMF: Network Embedding as Sparse Matrix Factorization

Experimental Results

# DeepWalk and NetMF



# DeepWalk and NetMF



Levy & Goldberg (NIPS 14)

$$\log \left( \frac{\#(w, c) |D|}{b \#(w) \#(c)} \right)$$

$b$  Number of negative samples

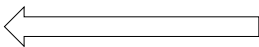
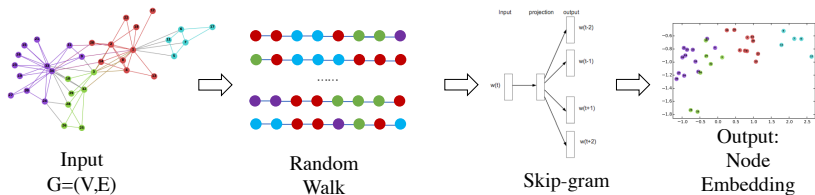
$\#(w, c)$  Co-occurrence of  $w$  and  $c$

$\#(w)$  Occurrence of word  $w$

$|D|$  Total number of word-context pairs

$\#(c)$  Occurrence of context  $c$

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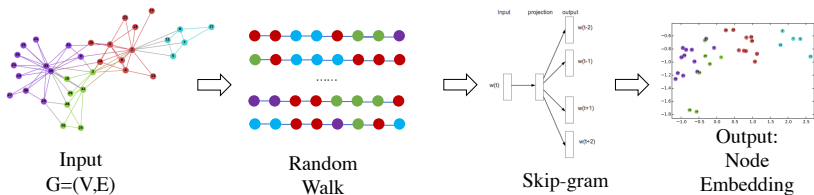
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# DeepWalk and NetMF

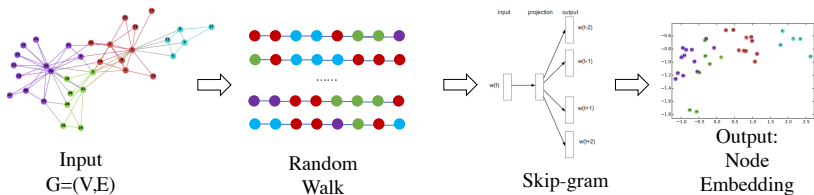


$$\log \left( \frac{\text{vol}(G)}{b} \left( \frac{1}{T} \sum_{r=1}^T (D^{-1}A)^r \right) D^{-1} \right)$$

$$\log \left( \frac{\#(w, c) |\mathcal{D}|}{b \#(w) \#(c)} \right)$$

- A** Adjacency matrix      $\text{vol}(G) = \sum_i \sum_j A_{i,j}$   
**D** Degree matrix          $b$  Number of negative samples

# DeepWalk and NetMF



$$\log \left( \frac{\text{vol}(G)}{b} \left( \frac{1}{T} \sum_{r=1}^T (D^{-1}A)^r \right) D^{-1} \right)$$

**Matrix  
Factorization**

$$\log \left( \frac{\#(w, c) |\mathcal{D}|}{b \#(w) \#(c)} \right)$$

Levy & Goldberg (NIPS 14)

Revisit DeepWalk and NetMF

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# Contents

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## Computation Challenges of NetMF

For small world networks,

$$\frac{\text{vol}(G)}{b} \left( \frac{1}{T} \sum_{r=1}^T \underbrace{(D^{-1}A)^r}_{\text{matrix power}} \right) D^{-1} \text{ is always a dense matrix .}$$

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Why?

- | In small world networks, each pair of vertices  $(i;j)$  can reach each other in a small number of hops.
- | Make the corresponding matrix entry a positive value.

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Idea

- | Sparse matrix is easier to handle.
- | Can we achieve a matrix sparse but 'good enough' matrix.

# Observation

## Definition

For  $\sum_{r=1}^T \alpha_r = 1$  and  $\alpha_r$  non-negative,

$$L = D \sum_{r=1}^T \alpha_r D^{-1} A^r D \quad (1)$$

is a  $T$ -degree random-walk matrix polynomial.

## Observation

For  $\alpha_1 = \dots = \alpha_T = \frac{1}{T}$ :

$$\begin{aligned}
& \log \left( \frac{\text{vol}(G)}{b} \left( \frac{1}{T} \sum_{r=1}^T (D^{-1} A)^r \right) D^{-1} \right) \\
&= \log \left( \frac{\text{vol}(G)}{b} D^{-1} (D \tilde{L} D^{-1}) \right) \\
&= \log \left( \frac{\text{vol}(G)}{b} D^{-1} (D \tilde{L} D^{-1}) \right)
\end{aligned}$$



# Random-walk Matrix Polynomial Sparsification

## Theorem

[CCL<sup>+</sup> 15] For random-walk matrix polynomial  $L = D \sum_{r=1}^T \alpha_r D^{-1} A^r$ , one can construct, in time  $O(T^2 m \log^2 n)$ , a  $(1 + \epsilon)$ -spectral sparsifier,  $\tilde{L}$ , with  $O(n \log n \log^2 n)$  non-zeros. For unweighted graphs, the time complexity can be reduced to  $O(T^2 m \log n)$ .

# NetSMF—Algorithm

The proposed NetSMF algorithm consists of three steps:

- | Construct a random walk matrix polynomial sparsifier,  $\tilde{L}$ , by calling PathSampling algorithm proposed in [CCL<sup>+</sup>15].
- | Construct a NetMF matrix sparsifier.

$$\text{trunc\_log} \left( \frac{\text{vol}(G)}{b} D^{-1} (D - \tilde{L}) D^{-1} \right)$$

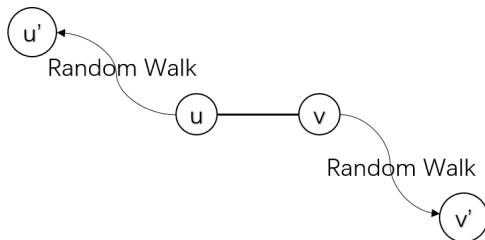
- | Truncated randomized singular value decomposition.

Detailed Algorithm

# Algorithm Details

## PathSampling:

- | Sample an edge  $(u; v)$  from edge set.
- | Start very short random walk from  $u$  and arrive  $u'$ .
- | Start very short random walk from  $v$  and arrive  $v'$ .
- | Record vertex pair  $(u'; v')$ .



## Randomized SVD:

- | Project origin matrix to low dimensional space by Gaussian random matrix.
- | Deal with the projected small matrix.

# NetSMF — System Design

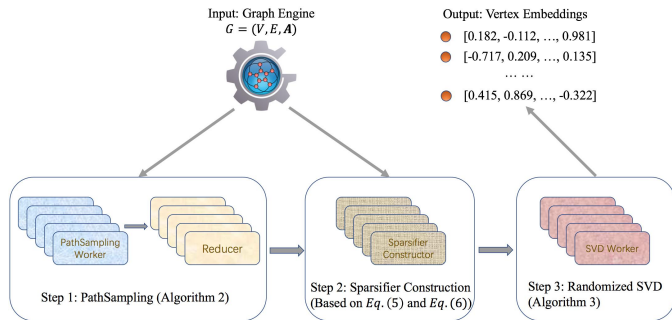


Figure 2: The System Design of NetSMF.

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# Setup

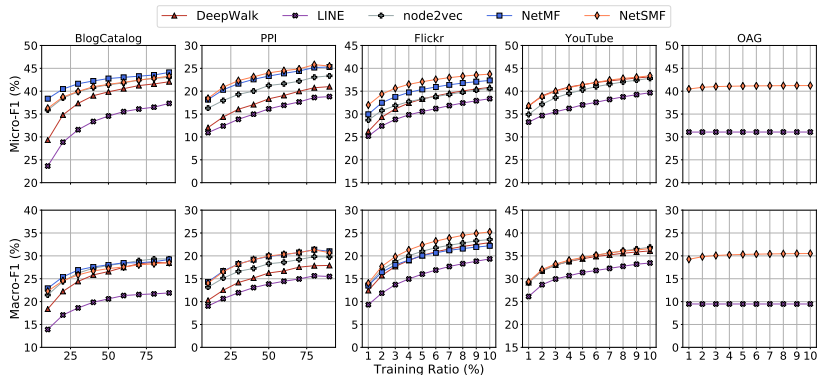
Label Classification:

- | BlogCatalog, PPI, Flickr, YouTube, OAG.
- | Logistic Regression
- | NetSMF ( $T = 10$ ), NetMF ( $T = 10$ ), DeepWalk, LINE.

Table 1: Statistics of Datasets.

Dataset	BlogCatalog	PPI	Flickr	YouTube	OAG
$jVj$	10,312	3,890	80,513	1,138,499	67,768,244
$jEj$	333,983	76,584	5,899,882	2,990,443	895,368,962
#Labels	39	50	195	47	19

# Experimental Results



**Figure 3:** Predictive performance on varying the ratio of training data. The x-axis represents the ratio of labeled data (%), and the y-axis in the top and bottom rows denote the Micro-F1 and Macro-F1 scores respectively.

# Running Time

Table 2: Running Time

	<b>LINE</b>	<b>DeepWalk</b>	<b>node2vec</b>	<b>NetMF</b>	<b>NetSMF</b>
BlogCatalog	40 mins	12 mins	56 mins	2 mins	13 mins
PPI	41 mins	4 mins	4 mins	16 secs	10 secs
Flickr	42 mins	2.2 hours	21 hours	2 hours	48 mins
YouTube	46 mins	1 day	4 days		4.1 hours
OAG	2.6 hours	–	–		24 hours



# Conclusion and Future Work

We propose NetSMF, a scalable, efficient, and effective network embedding algorithm.

## Future Work

- | A distributed-memory implementation.
- | Extension to directed, dynamic, heterogeneous graphs.

# Thanks.

- | Network Embedding as Matrix Factorization: Unifying DeepWalk, LINE, PTE, and node2vec (WSDM '18)
- | NetSMF: Network Embedding as Sparse Matrix Factorization (WebConf '19)

Code for NetMF available at [github.com/xptree/NetMF](https://github.com/xptree/NetMF)  
Code for NetSMF available at [github.com/xptree/NetSMF](https://github.com/xptree/NetSMF)

Q&A

# On the Large-dimensionality Assumption of [LG14]

Recall the objective of skip-gram model:

$$\min_{\mathbf{X}; \mathbf{Y}} L(\mathbf{X}; \mathbf{Y})$$

where

$$L(\mathbf{X}; \mathbf{Y}) = \sum_w \sum_c \left( \frac{\#(w; c)}{jDj} \log g(\mathbf{x}_w^> \mathbf{y}_c) + b \frac{\#(w)}{jDj} \frac{\#(c)}{jDj} \log g(\mathbf{x}_w^> \mathbf{y}_c) \right)$$

## Theorem

For DeepWalk, when the length of random walk  $L \gg 1$ ,

$$\frac{\#(w; c)}{jDj} \approx \frac{1}{2T} \sum_{r=1}^T \left( \frac{d_w}{\text{vol}(G)} (\mathbf{P}^r)_{w;c} + \frac{d_c}{\text{vol}(G)} (\mathbf{P}^r)_{c;w} \right);$$

$$\frac{\#(w)}{jDj} \approx \frac{d_w}{\text{vol}(G)} \text{ and } \frac{\#(c)}{jDj} \approx \frac{d_c}{\text{vol}(G)};$$

# NetSMF — Approximation Error

Denote  $\mathbf{M} = \mathbf{D}^{-1}(\mathbf{D} - \mathbf{L})\mathbf{D}^{-1}$  in

$$\text{trunc\_log} \left( \frac{\text{vol}(G)}{b} \mathbf{D}^{-1}(\mathbf{D} - \tilde{\mathbf{L}})\mathbf{D}^{-1} \right);$$

and  $\tilde{\mathbf{M}}$  to be its sparsifier the we constructed.

## Theorem

The singular value of  $\tilde{\mathbf{M}} - \mathbf{M}$  satisfies

$$\sigma_i(\tilde{\mathbf{M}} - \mathbf{M}) \leq \frac{4}{d_i d_{\min}}; \forall i \in [n];$$

## Theorem

Let  $\|\cdot\|_F$  be the matrix Frobenius norm. Then

$$\left\| \text{trunc\_log} \left( \frac{\text{vol}(G)}{b} \tilde{\mathbf{M}} \right) - \text{trunc\_log} \left( \frac{\text{vol}(G)}{b} \mathbf{M} \right) \right\|_F \leq \frac{4 \text{vol}(G)}{b d_{\min}} \sqrt{\sum_{i=1}^n \frac{1}{d_i}};$$

# Spectrally Similar

## Definition

Suppose  $G = (V; E; \mathbf{A})$  and  $\tilde{G} = (V; \tilde{E}; \tilde{\mathbf{A}})$  are two weighted undirected networks. Let  $\mathbf{L} = \mathbf{D}_G - \mathbf{A}$  and  $\tilde{\mathbf{L}} = \mathbf{D}_{\tilde{G}} - \tilde{\mathbf{A}}$  be their Laplacian matrices, respectively. We define  $G$  and  $\tilde{G}$  are  $(1 + \epsilon)$ -spectrally similar if

$$\forall \mathbf{x} \in \mathbb{R}^n; (1 - \epsilon) \mathbf{x}^T \tilde{\mathbf{L}} \mathbf{x} \leq \mathbf{x}^T \mathbf{L} \mathbf{x} \leq (1 + \epsilon) \mathbf{x}^T \tilde{\mathbf{L}} \mathbf{x}$$

# NetSMF—Algorithm

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## Algorithm 1: NetSMF

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**Input** : A social network  $G = (V; E; \mathbf{A})$  which we want to learn network embedding;  
The number of non-zeros  $M$  in the sparsifier; The dimension of embedding  $d$ .

**Output**: An embedding matrix of size  $n \times d$ , each row corresponding to a vertex.

```
1  $\mathcal{G} \leftarrow (V; ; \mathbf{A} = \mathbf{0})$ 
   /* Create an empty network with  $E = ;$  and  $\mathbf{A} = \mathbf{0}$ . */
2 for  $i = 1$  to  $M$  do
3     Uniformly pick an edge  $e = (u; v) \in E$ 
4     Uniformly pick an integer  $r \in [T]$ 
5      $u'; v'; Z \leftarrow \text{PathSampling}(e, r)$ 
6     Add an edge  $u'; v'; \frac{2rm}{MZ}$  to  $\mathcal{G}$ 
       /* Parallel edges will be merged into one edge, with their weights
       summed up together. */
7 end
8 Compute  $\mathbf{L}$  to be the unnormalized graph Laplacian of  $\mathcal{G}$ 
9 Compute  $\hat{\mathbf{M}} = \mathbf{D}^{-1} \mathbf{D} \mathbf{L} \mathbf{D}^{-1}$ 
10  $U_d; \Sigma_d; V_d \leftarrow \text{RandomizedSVD}(\text{trunc\_log} \circ \frac{\text{vol}(G)}{b} \hat{\mathbf{M}} ; d)$ 
11 return  $U_d \overset{p}{\Sigma_d}$  as network embeddings
```

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**Algorithm 2:** PathSampling algorithm as described in [CCL<sup>+</sup>15].

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- 1 **Procedure** PathSampling( $e = (u; v), r$ )
  - 2     Uniformly pick an integer  $k \in [r]$
  - 3     Perform  $(k - 1)$ -step random walk from  $u$  to  $u_0$
  - 4     Perform  $(r - k)$ -step random walk from  $v$  to  $u_r$
  - 5     Keep track of  $Z(\mathbf{p}) = \sum_{i=1}^r \frac{2}{A_{u_{i-1}, u_i}}$  along the length- $r$  path  $\mathbf{p}$   
      between  $u_0$  and  $u_r$
  - 6     **return**  $u_0; u_r; Z(\mathbf{p})$
- 

Back

# Randomized SVD

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## Algorithm 3: Randomized SVD on NetMF Matrix Sparsifier

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```
1 Procedure RandomizedSVD( $\mathbf{A}$ ,  $d$ )
2   Sampling Gaussian random matrix  $\mathbf{O}$  //  $\mathbf{O} \in \mathbb{R}^{n \times d}$ 
3   Compute sample matrix  $\mathbf{Y} = \mathbf{A}^\top \mathbf{O} = \mathbf{A}\mathbf{O}$  //  $\mathbf{Y} \in \mathbb{R}^{n \times d}$ 
4   Orthonormalize  $\mathbf{Y}$ 
5   Compute  $\mathbf{B} = \mathbf{A}\mathbf{Y}$  //  $\mathbf{B} \in \mathbb{R}^{n \times d}$ 
6   Sample another Gaussian random matrix  $\mathbf{P}$  //  $\mathbf{P} \in \mathbb{R}^{d \times d}$ 
7   Compute sample matrix of  $\mathbf{Z} = \mathbf{B}\mathbf{P}$  //  $\mathbf{Z} \in \mathbb{R}^{n \times d}$ 
8   Orthonormalize  $\mathbf{Z}$ 
9   Compute  $\mathbf{C} = \mathbf{Z}^\top \mathbf{B}$  //  $\mathbf{C} \in \mathbb{R}^{d \times d}$ 
10  Run Jacobi SVD on  $\mathbf{C} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^\top$ 
11  return  $\mathbf{Z}\mathbf{U}$ ,  $\mathbf{\Sigma}$ ,  $\mathbf{Y}\mathbf{V}$ 
/* Result matrices are of shape  $n \times d; d \times d; n \times d$  resp. */
```

---



# Time and Space Complexity

Table 3: Time and Space Complexity of NetSMF.

	Time	Space
Step 1	$O(MT \log n)$ for weighted networks $O(MT)$ for unweighted networks	$O(M + n + m)$
Step 2	$O(M)$	$O(M + n)$
Step 3	$O(Md + nd^2 + d^3)$	$O(M + nd)$

# References I

-  Dehua Cheng, Yu Cheng, Yan Liu, Richard Peng, and Shang-Hua Teng, *Spectral sparsification of random-walk matrix polynomials*, arXiv preprint arXiv:1502.03496 (2015).
-  Omer Levy and Yoav Goldberg, *Neural word embedding as implicit matrix factorization*, Advances in neural information processing systems, 2014, pp. 2177–2185.